

An Introduction To Gröbner Basis

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- While they mainly help in computations, they can also be used to prove substantial theorems, such as Hilbert Basis Theorem, Hilbert's Syzygy Theorem, etc. (Of course, any systematic study generates its own problems as well.)

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Given a finite set of generators for an ideal $I \subset k[X_1, \dots, X_n]$, can we find a finite set of generators for $I \cap k[X_{r+1}, \dots, X_n]$, $1 \leq r \leq n-1$.

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Check the Wikipedia page on Gröbner basis for many such computational problems where GB acts as a positive catalyst. Currently, it is being applied to applied fields like coding theory in error-correcting codes as well.

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A **monomial** is of the form $m = x_1^{\alpha_1} \dots x_n^{\alpha_n}$ ($:= \mathbf{x}^A$; $A = (\alpha_1, \dots, \alpha_n)$; Write $|A| = \sum_i \alpha_i$), where $\alpha_i \in \mathbb{Z}_+ \cup \{0\}$.

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A **term** is an element of the form λm where $\lambda \in k$, $m \in \mathcal{M}$.

Orderings

- A **term ordering** τ (denoted $>_{\tau}$) is a partial ordering on \mathcal{M} where:
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 - (b) if $a, b, c \in \mathcal{M}$ with $a >_{\tau} b$, then $ac >_{\tau} bc$.
- A **monomial ordering** τ (notated $>_{\tau}$) is a term ordering on \mathcal{M} such that $>_{\tau}$ is a total ordering.
- A **degree-wise monomial ordering** is a monomial ordering which respects the degrees of a monomial.

Example

Let $R = k[X, Y, Z]$ and τ be a monomial ordering such that $X >_{\tau} Y >_{\tau} Z$. Consider the degree 2 monomials.

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Thus, here we need to make a choice in degree 2, namely $XZ >_{\tau} Y^2$ or $Y^2 >_{\tau} XZ$.

Types of Monomial Orderings

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- 3 **Revlex:** $\mathbf{x}^A >_{revlex} \mathbf{x}^B$ if and only if either:
 - 1 $|A| > |B|$ or
 - 2 $|A| = |B|$ and the *last* nonzero entry of $A - B$ is negative.

Example

Let $R = k[x_1, x_2, x_3]$, $A = (4, 2, 6)$, and $B = (2, 3, 4)$.

Then:

- $A - B = (2, -1, 2)$

Therefore $\mathbf{x}^A >_{lex} \mathbf{x}^B$ since the first entry is positive.

Example

Let $R = k[x_1, x_2, x_3, x_4, x_5, x_6]$, $A = (4, 2, 6, 3, 1, 5)$, and $B = (4, 4, 0, 4, 4, 5)$.

Then:

- $|A| = 4 + 2 + 6 + 3 + 1 + 5 = 21$
- $|B| = 4 + 4 + 4 + 0 + 4 + 5 = 21$
- $A - B = (0, -2, 2, 3, -3, 0)$

So $|A| = |B|$.

Therefore $\mathbf{x}^B >_{deglex} \mathbf{x}^A$ since $\mathbf{x}^B >_{lex} \mathbf{x}^A$.

We also have that $\mathbf{x}^A >_{revlex} \mathbf{x}^B$ since the *last nonzero* entry is negative.

Initial Ideal

Let τ be a monomial ordering and let $f \in R = k[x_1, x_2, \dots, x_n]$.

- 1 The **initial term** of f with respect to τ (denoted $\text{in}_\tau(f)$) is the largest monomial in a term of f .

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- 1 The **initial term** of f with respect to τ (denoted $\text{in}_\tau(f)$) is the largest monomial in a term of f .
- 2 The **leading term** of f with respect to τ (notated $\text{lt}_\tau(f)$) is the term in f which has $\text{in}_\tau(f)$.
- 3 The **initial ideal** of I with respect to τ (denoted $\text{in}_\tau(I)$) is the ideal generated by the initial terms of all elements (**not necessarily just generators**) in I . Notationally, $\text{in}_\tau(I) := (\text{in}_\tau(f) : f \in I)$.

Example

Let $R = k[x, y]$ with $x >_r y$. Let f, g, h be polynomials in R where

$$f = x^3 + 3x^2y + 3xy^2 + y^3, g = 4x^2 - y^2, h = 3xy + 6y^2$$

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Example

$R = k[x, y, z]$, $\text{char}(k) \neq 2$, $I = (x + y - z, x - y + z)$, τ is a monomial ordering on R such that $x >_{\tau} y >_{\tau} z$. Let

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$$\exists f = \lambda z^i + (\text{lower terms only in } y \text{ and } z) \in I, \lambda \in k$$

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But, no such lower term can exist! So, $z^i \in I$ and hence $z \in \sqrt{I}$.

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- f **directly reduces** to h with respect to g if $\mu \mathbf{x}^A = \text{lt}_\tau(g)$ divides a nonzero term $\lambda \mathbf{x}^B$ of f and $h = f - \left(\frac{\lambda}{\mu}\right) \mathbf{x}^{B-A} g$.
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 - Notation: $f \xrightarrow{G} h$
- h is **reduced with respect to** G if no term in h is divisible by $\text{in}_{\tau}(g_i)$ for any $g_i \in G$

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$$f = \sum_{i=1}^k u_i g_i + r \quad (*)$$

where r is reduced with respect to G and

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$$\text{in}_\tau(f) \geq \max\{\text{in}_\tau(u_1 g_1), \dots, \text{in}_\tau(u_k g_k)\}$$

r : remainder w.r.t. G (is it unique?)

The Division Algorithm

Theorem

Let $G = \{g_1, \dots, g_k\}$ be a collection of nonzero polynomials in R and f be any polynomial in R . There are polynomials $u_1, \dots, u_k, r \in R$ such that we can write

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where r is reduced with respect to G and

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Notation: $r = \bar{f}^G$.

Gröbner Basis

Definition

A Gröbner basis of an ideal I in the polynomial ring $k[x_1, x_2, \dots, x_n]$ with respect to the monomial ordering τ is a set $\{f_1, \dots, f_m\} \subset I$ such that

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- Recall that we found $\text{in}_\tau(I) = (x, y)$ for $I = (x - y + z, x + y - z)$. Hence a Gröbner Basis is given $\{x + y - z, 2y - 2z\}$.

Gröbner Basis Generates I

Theorem

Let J, I be ideals of R such that $J \subset I$. Let τ be a monomial ordering on R . Then

$$\text{in}_\tau(I) = \text{in}_\tau(J) \iff I = J.$$

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Corollary

If $G = \{g_1, \dots, g_r\}$ is a Gröbner basis of the ideal I , then G generates I .

Criterion for Ideal Membership

Theorem

Let $I \subseteq R$ be an ideal and let $G = \{g_1, \dots, g_k\} \subseteq I$, $g_i \neq 0$ for all i . Then the following are equivalent:

- 1 G is a Gröbner basis for I .
- 2 $f \in I$ if and only if $f \xrightarrow{G} 0$.

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Let $R = k[x, y, z]$. τ : deglex with $x >_\tau y >_\tau z$. Let $I = (g_1, g_2)$ where $g_1 = xy - z^2$ and $g_2 = y^2 - xz$.

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Theorem (Buchberger's Criterion)

Let τ be a monomial ordering on R and $I = (g_1, \dots, g_s)$ be an ideal. Let $G = \{g_1, \dots, g_s\}$. Fix remainders of $S(g_i, g_j)$ with respect to G , say $\overline{S(g_i, g_j)}^G$. Then $G = \{g_1, \dots, g_s\}$ is a Gröbner basis for I if and only if $\overline{S(g_i, g_j)}^G = 0$ for all i, j where $1 \leq i < j \leq s$.

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This process stops.

Reason: We get an ascending chain

$$(\text{in}_\tau(g_1), \dots, \text{in}_\tau(g_s)) \subseteq (\text{in}_\tau(g_1), \dots, \text{in}_\tau(g_s), \text{in}_\tau(h_{ij})) \subseteq \dots$$

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in the Noetherian ring R , hence stabilizes i.e. the S -polynomials reduce to 0 after finitely many steps.

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Now we start the process over again adjoining $S(g_1, g_2) = y^3 - z^3$ to G .

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- 1 Any set of monomials is a Gröbner basis for the ideals they generate.

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- ⑤ If $\{g_1, \dots, g_s\}$ is a Gröbner basis for $I = (g_1, \dots, g_s)$ and $k \subseteq L$ is a field extension, then $\{g_1, \dots, g_s\}$ is a Gröbner basis for $IL[x_1, \dots, x_n]$.

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- 2 Let $\{f_1, \dots, f_s\}$ be a Gröbner basis of I . Then

$$\{f_1, \dots, f_s\} \cap k[x_i, x_{i+1}, \dots, x_n]$$

is a Gröbner basis of $I \cap k[x_i, x_{i+1}, \dots, x_n]$.

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